| FIRST:   |                   |        |                      |             | LA             | ST:   | :               |      |    |           |  |
|--|-------------------|--------|----------------------|-------------|----------------|-------|-----------------|------|----|-----------|--|
| Discussion:  | Monday            | 9      | 10                   | 11          | 12             | 2     | 1               | 2    | 3  | 4         | 5  |
| (15 points) Check  | k the (single) bo | ox tha | at best              | chara       | acteri         | zes e | each            | iten | n. |           |  |
| T(1) = d $T(n) = 4T(n/2) - 4T(n/2$ | + n               | Θ(     | $n)$ $n^{\log_3 2})$ |             | ]              | ,     | $\log \log_2 3$ | ŕ    |    |           | $\Theta(n^2)$ $$ $\Theta(2^n)$ $\square$ |
| Producing all part for a sentence.   | rses              | po     | lynomi               | al [        |                | ex    | kpon            | enti | al | $\sqrt{}$ | in NP                                    |
| The running timesively defined by  |                   |        |                      |             | (n-1) $T(n/2)$ | ,     | _               |      |    |           | $(-1) + cn$ $(n/2) + cn$ $\sqrt{}$       |
| T(1) = d   |                   | )      |                      | $\chi(m^2)$ |                | 7 0   | V ~ 1           | ,    | \  | 7 (       | $\mathcal{N}(2^n)$                       |

$$T(n) = T(n-1) + n \qquad \qquad \Theta(n) \qquad \boxed{ } \qquad \Theta(n^2) \qquad \boxed{ } \qquad \Theta(n \log n) \qquad \boxed{ } \qquad \Theta(2^n) \qquad \boxed{ } \qquad \boxed{ }$$

The running time of Karatsuba's algorithm is recursively defined by T(1)=d and T(n)=

$$2T(n/2) + cn \qquad 3T(n/2) + cn \qquad \boxed{\checkmark}$$

$$4T(n/2) + cn$$
  $4T(n/2) + c$ 

| CS 173, Spring 2016 Examlet 11, Part B NETID:  |
|--|
| FIRST: LAST:   |
| Discussion: Monday 9 10 11 12 1 2 3 4 5  |
| (15 points) Check the (single) box that best characterizes each item.  |
| $\Theta(n) \qquad \Theta(n \log n) \qquad \sqrt{}$ The running time of mergesort $\Theta(n^2) \qquad \Theta(2^n) \qquad \boxed{}$  |
| $T(1) = d$ $T(n) = 2T(n/2) + c$ $\Theta(\log n)  \square$ $\Theta(n)   \qquad \Theta(n \log n)  \square$ $\Theta(n^2)  \square$  |
| For a problem to satisfy the definition of NP, a "no" answer must have a succinct justification. true $\ $ false $\ $  |
| This question had a typo of "x+2" for "x+1". We answered individual questions at the exam rat an posting a general correction because the latter would have given the question away. |
| Algorithm A takes $\log_2 n$ time. On one input, A takes x time. How long will it take if I double the input size? $x+1$ $\sqrt{}$ $2x$ $2x$ $x^2$ $x^2$                             |
| The chromatic number of a graph with $n$ nodes can be found in polynomial time. true false not known $\sqrt{}$   |

| CS 173, Spring 2016 Examlet 11, Part B NETID:   |
|---|
| FIRST: LAST:  |
| Discussion: Monday 9 10 11 12 1 2 3 4 5   |
| (15 points) Check the (single) box that best characterizes each item.   |
| The running time of the Towers of Hanoi solver is recursively defined by $T(1)=d$ and $T(n)=$ $2T(n-1)+c  $ $2T(n-1)+c n  $ $2T(n/2)+c  $ $2T(n/2)+c n  $ |
| For a problem to satisfy the definition of co-NP, a "yes" answer must have a succinct justification. true $\ $ false $\ $                                 |
| The running time of the Towers $\Theta(\log n)$ $\Theta(n\log n)$ $\Theta(n\log n)$ of Hanoi solver $\Theta(n^2)$ $\Theta(2^n)$ $O(2^n)$                  |
| $n^{\log_2 3}$ grows faster than $n$ slower than $n$ at the same rate as $n$  |
| The Marker Making problem can be solved in polynomial true false not known $\sqrt{}$  |

| Examlet 11, Pa  |                                |                    | LAST                        | ף•                |         |       |         |           |      |
|---|--------------------------------|--------------------|-----------------------------|-------------------|---------|-------|---------|-----------|------|
| rittsi.   |                                |                    |                             | L •               |         |       |         |           |      |
| Discussion: Mo  | nday 9                         | 10 11              | 12                          | 1 2               | 2 3     | 4     | 5       |           |      |
| (6 points) Your partners of integers. Using Merg  | -                              |                    |                             |                   | , .     |       | _       |           | link |
| $Mergesort(L = (a_1, a_2))$   | $,\ldots,a_n)) \setminus$      | input is a         | linked lis                  | t L con           | taining | n int | egers   |           |      |
| Solution: if (n   | =1) return L                   |                    |                             |                   |         |       |         |           |      |
|   |                                |                    |                             |                   |         |       |         |           |      |
| p = floor(n/2)  |                                |                    |                             |                   |         |       |         |           |      |
| Solution:   |                                |                    |                             |                   |         |       |         | 7         |      |
| $L_a=(a_1,\ldots$   | • .                            |                    |                             |                   |         |       |         |           |      |
| $L_b = (a_{p+1}, $ return Mer   | $\ldots, a_n$ ) ge(Mergesort(  | $L_a$ ), Merge     | $\operatorname{esort}(L_b)$ |                   |         |       |         |           |      |
|   |                                |                    |                             |                   |         |       |         |           |      |
| (9 points) Check the (six   | ngle) box that                 | best chara         | acterizes e                 | each ite          | m.      |       |         |           |      |
| T(1) = d<br>T(n) = 2T(n-1) + c  | $\Theta(n)$                    | $\Theta(n^2)$      | ) [ (                       | $\Theta(n\log n)$ | n)      | Θ     | $(2^n)$ | $\sqrt{}$ |      |
| It takes exponential time<br>a propositional logic ex-<br>true by picking the rig<br>its propositional variable | xpression can<br>ht true/false | be made values for | true                        |                   | false   |       | ]       | not known |      |
| its propositional variable  | es (e.g. p, q, r               | ·).                | _                           |                   |         |       |         |           |      |

| CS 173, Spring 2016 Examlet 11, Part B NETID:  |  |  |  |  |  |
|--|--|--|--|--|--|
| FIRST: LAST:   |  |  |  |  |  |
| Discussion: Monday 9 10 11 12 1 2 3 4 5  |  |  |  |  |  |
| (15 points) Check the (single) box that best characterizes each item.  |  |  |  |  |  |
| Circuit satisfiability can be solved in exponential time. true $\sqrt{}$ false $\overline{}$ not known $\overline{}$   |  |  |  |  |  |
| $\Theta(n^2) \qquad \Theta(n^3) \qquad \Theta(n\log n) \qquad \\ \text{Karatsuba's integer} \\ \text{multiplication algorithm} \qquad \Theta(n^{\log_2 3}) \qquad \boxed{\checkmark} \qquad \Theta(n^{\log_3 2}) \qquad  \qquad \Theta(2^n) \qquad $ |  |  |  |  |  |
| Marker Making polynomial exponential $\sqrt{}$ in NP   |  |  |  |  |  |
| The running time of mergesort is $\Theta(n^3)$ . true false $\sqrt{}$  |  |  |  |  |  |
| $n^{\log_2 4}$ grows faster than $n^2$ slower than $n^2$ at the same rate as $n^2$   |  |  |  |  |  |

| CS 173, Spring 20<br>Examlet 11, Part  | NETID:   |
|--|--|
| FIRST:   | LAST:  |
| Discussion: Monda  | y 9 10 11 12 1 2 3 4 5   |
| (15 points) Check the (single  | ) box that best characterizes each item.   |
| The Travelling Salesman  | polynomial exponential in NP $\sqrt{}$   |
| Merging two sorted lists of numbers $\Theta(\log n)$                                       | $\Theta(n)$ $\Theta(n)$ $\Theta(n \log n)$ $\Theta(n^2)$                                     |
| The running time of the Towe of Hanoi solver   | $\Theta(\log n)$ $\Theta(n \log n)$ $\Theta(n \log n)$ $\Theta(n^2)$ $\Theta(2^n)$ $O(2^n)$  |
| T(1) = d $T(n) = 2T(n/4) + c$  | $\Theta(\sqrt{n})$ $$ $\Theta(n)$ $$ $\Theta(n \log n)$ $$ $\Theta(n^2)$ $$ $\Theta(2^n)$ $$ |
| Algorithm A takes $2^n$ time. input, A takes x time. How left take if I double the input s | ong will $x + 2$ $2x$ $2x$ $2^x$ $x^2$   |