University of Illinois at Urbana-Champaign First Midterm Exam, ECE 220 Honors Section

Thursday 15 February 2018

	Nam	e:					
	Net I	D:					
•	Be su	re that your ex	xam booklet has ELEVEN pages.				
•	Write your name and Net ID on the first page.						
•	Do not tear the exam apart other than to remove the reference sheet.						
•	• This is a closed book exam. You may <u>not</u> use a calculator.						
•	You are allowed one handwritten 8.5×11-inch sheet of notes (both sides).						
•	 The last page of the exam gives RTL for LC-3 instructions (except JSRR; given JSRR BaseR, the RTL is PC ← BaseR, R7 ← PC). Copies of Patt & Patel's Appendix A are also available during the exam. Absolutely no interaction between students is allowed. 						
•							
•	Show	all work, and	clearly indicate any assumptions that you make.				
•	Don't	t panic, and go	od luck!				
Proble	em 1	24 points					
Proble	em 2	20 points					
Proble	em 3	20 points					
Proble	em 4	16 points					
Proble	em 5	20 points					
Total		100 points					

Problem 1 (24 points): Short Answer Questions

1. (12 points) A fellow ECE 220 student is attempting to write a program that calculates either the product or the sum of numbers from 1 to N. The user first types the value of N, which must be a positive integer. The student's code reads the number N from the keyboard using the READ N subroutine, which handles any non-numeric input as well as overflow, returning only when the human user has entered a positive value for N. Next, the user presses either the asterisk key ("*") or the plus key ("+") to tell the program whether to me

calcula memor	te the product or the sum of the numbers from 1 to N, respectively. The result is then stored in y.
Using I credit.	NO MORE THAN 15 WORDS, describe each of the following. Answering with code will earn no
a.	(4 points) One subtask for which the programmer should use a sequential decomposition.
b.	(4 points) One subtask for which the programmer should use a conditional decomposition.
c.	(4 points) One subtask for which the programmer should use an iterative decomposition.
charact	points) Fill in the blanks below so that the resulting LC-3 subroutine, MYOUTPUT, writes an ASCII ter from R5 out to the display using memory-mapped I/O (in other words, do not use TRAPs nor other tine calls). You may change R0 and/or R7 if necessary. All other registers are callee-saved.
LUOYW	PUT
	RET
	VE I

DSR .FILL xFE04 DDR .FILL xFE06

Problem 1, continued:

3. **(8 points)** Consider the multiplication subroutine shown below.

; Tak	es two positive integers and returns the product							
MULT	AND R1, R1, #0							
	ADD R3, R3, #0							
T 00D	BRZ DONE							
LOOP	ADD R1, R1, R2							
	ADD R3, R3, #-1 BRp LOOP							
DONE	RET							
20112								
a.	a. (2 points) Which registers hold the operands for multiplication by the subroutine?							
_								
b. (2 points) In which register does the subroutine return the product?								
c	(2 points) Which registers in the MULT subroutine are callee-saved?							
	(2 points) which registers in the MOL1 subroutile are cance-saved:							
d.	(2 points) Which registers in the MULT subroutine are caller-saved?							

Problem 2 (20 points): Writing LC-3 Code

In this problem, you must write a subroutine, SWAPBYTES, that swaps the two 8-bit bytes held in one 16-bit LC-3 register.

For example, given the initial value xCDAB, your subroutine must move xCD into the low 8 bits and xAB into the high 8 bits to produce the value xABCD. Similarly, given the initial value xEBEC, your subroutine must swap the xEB with the xEC to produce the 16-bit value xECEB.

The call interface for SWAPBYTES is as follows:

Input: **R1**, a 16-bit value (2 bytes)

Output: R2, a 16-bit value (2 bytes) with bytes swapped

All registers are caller-saved.

Requirements for your subroutine:

• Use at most 25 instructions (excludes labels, .ORIG, .END,. FILL, .BLKW).

Any code after the 25th instruction will not be graded.

It is possible to finish this problem with 12 instructions or less.

- You must use a loop(s) in your code. Manual repetition will receive **ZERO** credit.
- Briefly comment your code and describe how each register is used. You are not required to comment every line.

Write your subroutine on the next page.

(Use the back of this page if you need more space.)

Problem 2, continued: (call interface and examples reproduced for your convenience)

```
SWAPBYTES
Input: R1 - 2 bytes (16 bits)
Output: R2 - 2 bytes (16 bits) with bytes swapped
All registers are caller-saved.
Example 1:
                                            Example 2:
R1 - 0xCDAB → R2 - 0xABCD
                                            R1 - 0xEBEC → R2 - 0xECEB
 .ORIG x3100
 SWAPBYTES
 .END
```

Problem 3 (20 points): Division with a Stack

Write subroutine STACKDIVIDE in the space below. Given a stack containing some number (≥ 2) of positive integers, with R6 pointing to the top of the stack, the routine STACKDIVIDE operates as follows:

- 1. Pops two numbers.
- 2. Uses the DIVIDE subroutine (interface specified below) to divide the first number popped by the second number popped, ignoring any remainder.
- 3. Jumps to **Step 6** if the stack is empty (stack pointer equal to value at BASEPTR).
- 4. Otherwise, pushes the result onto the stack
- 5. Returns to **Step 1** (repeats the process).
- 6. Stores the result of the final division into the memory location labeled ANS.

Step 2 must utilize the DIVIDE subroutine, whose address is stored at the label DIVIDE_ADDRESS. As input, the DIVIDE subroutine expects the dividend (numerator) in R1 and the divisor (denominator) in R2, then returns the quotient in R3. All registers except R3 and R7 are callee-saved with DIVIDE.

On return from STACKDIVIDE, R6 must point to the base of the stack. All other registers are caller-saved. Use the memory location labeled as SAVE to save any value that you find necessary to save. Briefly comment your code and describe how each register is used. Use the back of <u>this</u> sheet if you need more space, but we will only read up to 30 instructions.

STACKDIVIDE

RET

BASEPTR .FILL x4000; base of stack

SAVE .BLKW #1 ANS .BLKW #1

DIVIDE ADDRESS .FILL x7000; the DIVIDE subroutine is at this address

Problem 4 (16 points): C Variables and Function Calls

1. (12 points) Read the program below.

```
#include <stdint.h>
#include <stdio.h>
int32 t mystery (int32 t x);
int main()
    int32 t a;
    int32 t b = 3;
    int32_t c = 5;
    int32 t d = 7;
    {
        int32 t c = 9;
        a = mystery (b);
        d = 11;
        printf ("a: %d, b: %d, c: %d, d: %d\n", a, b, c, d);
    }
   b = 7;
    a = mystery (b);
   printf ("a: %d, b: %d, c: %d, d: %d\n", a, b, c, d);
}
int32 t mystery (int32 t x)
   static int32 t y = 0;
    x = x + 1;
    y = y + x;
    return y;
}
```

Circle EXACTLY ONE ANSWER for each question.

- a. (6 points) For the two calls to printf(), what are the expected printed values of a and b?
 - 1) a: 4, b: 3 and a: 8, b: 7
 - 2) a: 4, b: 3 and a: 12, b: 7
 - 3) a: 4, b: 4 and a: 8, b: 8
 - 4) a: 4, b: 4 and a: 12, b: 8
 - 5) The program does not compile.
- b. (6 points) For the two calls to printf(), what are the expected printed values of c and d?
 - 1) c: 5, d: 7 and c: 5, d: 7
 - 2) c: 5, d: 11 and c: 5, d: 11
 - 3) c: 9, d: 11 and c: 5, d: 11
 - 4) c: 9, d: 11 and c: 9, d: 7
 - 5) The program does not compile.

Problem 4, continued:

2. (4 points) Consider a C function **example** called by another function. The left block of code below corresponds to the caller, and the right block corresponds to the callee (the function **example**).

This question focuses on whether the LC-3 instructions **depend on the number of parameters** needed by the function **example**.

For each of the four blanks (explained by the comments immediately above them), write "YES" if the LC-3 instructions depend on the number of parameters passed to **example**, or write "NO" if the LC-3 instructions do not depend on the number of parameters passed to **example**.

; prepare for call	EXAMPLE ; set up stack frame
JSR EXAMPLE	; (execute C statements)
; clean up after call	; tear down stack frame

RET

Problem 5 (20 points): Understanding Compiled C Code

1. **(15 points)** The LC-3 code below corresponds to the output of a compiler for the C function **foo**. Based on the LC-3 code, write C code for **foo** from which a non-optimizing compiler might have produced the LC-3 code.

```
FOO
     ADD R6, R6, \#-4
     STR R5, R6, #1
     ADD R5, R6, #0
     STR R7, R5, #2
     LDR R0, R5, #4
     STR R0, R5, #0
LOOP LDR R0, R5, #4
     BRnz DONE
     LDR R1, R5, #5
     ADD R6, R6, #-1
     STR R1, R6, #0
     LDR R2, R5, #0
     ADD R6, R6, \#-1
     STR R2, R6, #0
     JSR ANOTHER
     LDR R3, R6, #0
     ADD R6, R6, #3
                            ; call this subroutine "another" in C
     STR R3, R5, #0
     LDR R1, R5, #4
     ADD R1, R1, #-1
     STR R1, R5, #4
     BRnzp LOOP
DONE LDR RO, R5, #0
     STR R0, R5, #3
     LDR R7, R5, #2
     LDR R5, R5, #1
     ADD R6, R6, #3
     RET
```

Write the C function **foo** below. For parameters, choose names from X, Y, and Z. For local variables, choose names from A, B, and C. (There are no more than three of either type.) All types are **int**.

Problem 5, continued:

2. **(5 points)** Given the LC-3 implementation shown below of the C function **another**, write an expression for the value returned from function **foo** (from **Part 1**) in terms of the arguments passed (called **A** and **B** in the expression below):

foo (A, B) evaluates to ______.

ANOTHER

```
ADD R6, R6, #-3
STR R5, R6, #0
ADD R5, R6, #-1
STR R7, R5, #2
LDR R1, R5, #4
LDR R2, R5, #5
NOT R2, R2
ADD R2, R2, #1
ADD R1, R1, R2
STR R1, R5, #3
LDR R7, R5, #2
LDR R5, R5, #1
ADD R6, R6, #2
RET
```