ECE 220: Computer Systems & Programming

Spring 2019 - Midterm Exam 2

April 4, 2019

1.	This is a closed book exam except for 1 sheet of hand-written
	notes

- 2. You may not use any personal electronic devices, such as cellphone and calculator
- 3. Absolutely no interaction between students is allowed
- 4. Illegible handwriting will be graded as incorrect

Select the Lecture Section You Will Attend to Pick Up Your Booklet

[] Prof. Moon] Prof. Hu] Prof. Chen] Prof. Bhowmik	(BL2, 11:00 (BL, 12:30)	am) pm)
Nai	ne:		
Net	ID:		_
Roc	om:		_
Question 1 (20 p	ooints):		
Question 2 (20 p	oints):		
Question 3 (35 p	oints): A; B	; C	_
Question 4 (25 p	oints): A; B	; C	_; D
Total	Score:		

Problem 1 (20 points): Spiral Matrix

In this problem, we want to get the elements of an $m \times n$ matrix in **clockwise** spiral order. We will start the spiral from [0, 0].

$$\begin{bmatrix}
1 & 2 & 3 \\
4 & 5 & 6 \\
7 & 8 & 9
\end{bmatrix}
\rightarrow
\begin{bmatrix}
1 & 2 & 3 & 6 & 9 & 8 & 7 & 4 & 5
\end{bmatrix}$$

$$\begin{bmatrix}
1 & 2 & 3 \\
4 & 5 & 6 \\
7 & 8 & 9 \\
10 & 11 & 12
\end{bmatrix}
\rightarrow
\begin{bmatrix}
1 & 2 & 3 & 6 & 9 & 12 & 11 & 10 & 7 & 4 & 5 & 8
\end{bmatrix}$$

One approach to solving this problem is to unravel each outer layer of the matrix until the center is reached.

Below is the algorithm for solving this layer-by-layer:

- For each outer layer, iterate through the elements clockwise starting from the **top-left corner**:
 - o Traverse the top row of the layer
 - o Traverse the rightmost column of the layer
 - o If there are 4 sides to the layer
 - Traverse the bottom row of the layer
 - Traverse the leftmost column of the layer
 - Update the outer layer boundaries

Complete the following function with the following function declaration:

int* spiralMatrix(int* matrix, int RowSize, int ColSize);

- int* matrix → The 1D representation of the input 2D matrix (as seen on the left in the examples) that will be converted
- int RowSize → Number of rows in the input matrix
- int ColSize → Number of columns in the input matrix
- Returns 1D matrix containing the clockwise spiral

```
int* spiralOrder(int* matrix, int RowSize, int ColSize) {
   /* Allocate space for spiral matrix */
   int* ans = (int*) malloc(sizeof(int) * RowSize * ColSize);
   /* Set bounds for the outermost layer such that
     r low is the first row */
   int r low = 0;
   int r high = RowSize - 1;
   int c_{low} = 0;
   int c high = ColSize - 1;
   int c, r; /* Column and row indices */
   int idx = 0; /* Index for spiral matrix */
   // Iterate while there are outer layers
   while (r low<=r high && c low<=c high){
      /* Traverse right/down along row/column and
         add element to spiral matrix */
      for (c = c_low;_____; c++, idx++){
          ans[idx] = _____;
      for (r = r_low + 1;_____; r++, idx++){
          ans[idx] = _____;
       }
      /* Determine if it is possible to move left or up in
         this layer */
      if (r low<=r high && c low<=c high){</pre>
          /* Traverse left/up along row/column and
            add element to spiral matrix */
          for (c = c_high - 1;_____; c--, idx++){
             ans[idx] = _____;
          for (r = r_high;_____; r--, idx++){
             ans[idx] = _____;
          }
      }
      /* Update bounds for iteration (row low, row high,
                                  col low, col high) */
      r low++;
      r high--;
      c low++;
      c high--;
   return ans;
}
```

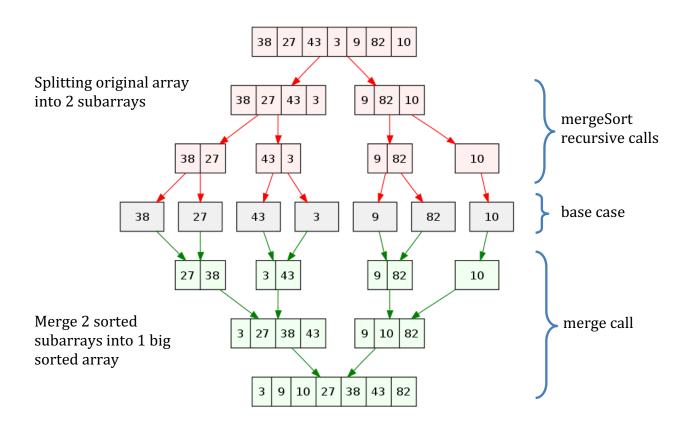
Problem 2 (20 points): Recursion - Merge Sort

In this problem, you need to implement an algorithm called 'Merge Sort' to sort the integer elements stored in an array in <u>ascending order</u>.

Input: 38, 27, 43, 3, 9, 82, 10

Output: 3, 9, 10, 27, 38, 43, 82

A schematic diagram of how merge sort works can be found below.



We will be using two functions – **merge** and **mergeSort** to implement this algorithm which are declared as following.

```
void merge(int arr[], int 1, int m, int r);
void mergeSort(int arr[], int 1, int r);
```

Now in this question, you need to implement the mergeSort function. Assume that the merge function has been implemented for you. **Hint: you need to call merge in mergeSort. Think about what should be the base case and the recursive case(s) in mergeSort.**

```
void merge(int arr[], int l, int m, int r){
    /* Implementation omitted for simplicity */
    /* It will merge two subarrays of arr[].
       First subarray is arr[l..m]
       Second subarray is arr[m+1..r] */
}
/* Use no more than 6 semicolons (;) in mergeSort, or you will
receive 0 for this question */
void mergeSort(int arr[], int 1, int r){
    /* l is the lowest index and r is the highest index
       of the arr to be sorted */
    /* base case */
    /* recursive case(s) */
}
int main(){
    int arr[] = {38, 27, 43, 3, 9, 82, 10};
    int arr size = 7;
    printf("Printing unsorted array\n");
    printArray(arr, arr_size); /* assume printArray given */
    mergeSort(arr, 0, arr_size - 1);
    printf("Printing sorted array\n");
    printArray(arr, arr size); /* assume printArray given */
    return 0;
}
```

Problem 3 (35 points): C to LC-3 Conversion

Tower of Hanoi is a mathematical puzzle where we have three rods and n disks. The objective of the puzzle is to move the entire stack to another rod, obeying the following simple rules:

- 1) Only one disk can be moved at a time.
- 2) Each move consists of taking the upper disk from one of the stacks and placing it on top of another stack i.e. a disk can only be moved if it is the uppermost disk on a stack.
- 3) No disk may be placed on top of a smaller disk.

```
Take an example for 2 disks: Let from_rod = 'A', mid_rod = 'B', to_rod = 'C'.

Step 1: Shift first disk from 'A' to 'B'.

Step 2: Shift second disk from 'A' to 'C'.

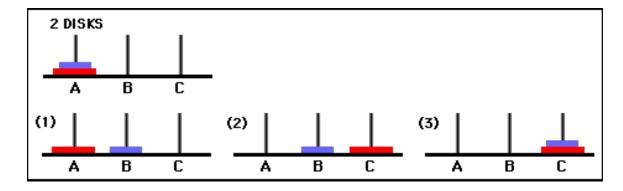
Step 3: Shift first disk from 'B' to 'C'.

The pattern here is:

Shift 'n-1' disks from 'A' to 'B'.

Shift last disk from 'A' to 'C'.

Shift 'n-1' disks from 'B' to 'C'.
```



```
// C recursive function to solve tower of hanoi puzzle
void towerOfHanoi(int n, char from_rod, char to_rod, char mid_rod) {
   if (n == 1) {
       printf("\n Move disk 1 from rod %c to rod %c", from rod,
                 to rod);
        return;
   towerOfHanoi(n-1, from rod, mid rod, to rod);
   printf("\n Move disk %d from rod %c to rod %c", n, from rod,
           to_rod);
   towerOfHanoi(n-1, mid rod, to rod, from rod);
}
int main() {
    int n = 3; // Number of disks
   towerOfHanoi(n, 'A', 'C', 'B'); // A, B and C are names of rods
   return 0;
}
```

Part A (3 points): For n = 3, what is the first three output? P	lease fill in the blanks:
Move disk from rod to rod	
Move disk from rod to rod	_
Move disk from rod to rod	<u> </u>
Part B (5 points) You MUST use and conform to the LC-3 caclass . Let R6 be the stack pointer and R5 to of stack when the first line of code (e.g. if (Use arrows to mark where R5 and R6 are	the frame pointer. Please show the status n==1)) inside towerofHanoi is executed.

Main's Activation Record

Part C (17 points): C to LC-3 Conversion

You **MUST** use and conform to the **LC-3 calling conventions** we have described **in class**. You do not have to convert the section denoted as "Omitted for simplicity". To receive full credit, you **MUST** complete each task within the number of lines provided in each section. A section is defined as a set of blank lines without any intervening provided code. A line is defined as a single instruction in LC-3. If you use a register that was set in a previous task without re-loading that register from the stack, you **MUST** document (with a comment) your usage of that register on first use within each section. Semicolons setting off space for comments have been provided on the given lines.

Please finish the LC-3 subroutine for the towerOfHanoi. You must use R2 to represent n, R3 for from_rod, R4 for mid_rod, R5 for to_rod.

towerOfHanoi ; Complete callee stack build up within 7 lines of code ; You MUST comment on each line of your code ; preparation for function logic _____; get the number of disks _____; get the from rod _____; get the to rod _____; get the mid rod ; Test for the termination condition (n == 1)AND R1, R1, #0 ADD R1, R1, #-1 ADD R1, R1, R2 BRz DisplayAndReturn ADD R2, R2, #-1 ; Update R2 on the stack;

<pre>; caller built-up for first recursive call ; omitted for simplicity ; H(n-1, source, dest, helper) JSR towerOfHanoi</pre>	
; Pop arguments off the stack	
;	
JSR Display ; caller built-up for second recursive call ; omitted for simplicity JSR towerOfHanoi BRnzp RETURN	
DisplayAndReturn ; Omitted for simplicity DisplayAndReturn ; Omitted for simplicity	
RETURN ; Complete callee stack tear-down within 6 lines of code ; You MUST comment on each line of your code	le
;	
;	
;	
;	
;	
;	
RET	

Problem 4 (25 points): Concepts

<pre>Part A (10 points): #include<stdio.h></stdio.h></pre>
<pre>int sub1(int arg); int t = 1;</pre>
int main (){
<pre>int sub1(int fluff){ int i = t; return (fluff + i); }</pre>
What is the output of the program?
First printf statement:
Second printf statement:

Part B (5 points):	
#include <stdio.h></stdio.h>	
<pre>typedef struct recordStruct{</pre>	
char name[100];	
char phone_number[10];	
}record;	
int main(){	
	Jane Doe", "123456789"};
record *ptr = &new_re	
ptr++;	,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,
printf("%c", &ptr);	
}	
What is the output of the pro-	gram?
r	5
Your answer:	
Part C: True or False (5 point	
	nother function, the callee function will have a local
copy of all the elements of the	e array.
W	
Your answer:	
Part D (5 points):	
	or in MP - 2048. What does it do?
•	f a structure directly through a variable
B. Have the same function	, e
	re pointer and access the member
D. Add more members to	•
Your answer:	(select all that apply)

Table E.2 The Standard ASCII Table ASCII ASCII ASCII ASCII Character Dec Hex Character Dec Hex Character Dec Hex Character Dec Hex 00 32 20 64 40 96 60 nul 0 sp 0 97 soh 1 01 33 21 65 41 61 Α stx 2 02 34 22 66 42 98 62 В b 3 43 99 etx 03 # 35 23 C 67 C 63 4 36 04 24 D 68 44 d 100 64 eot \$ 5 05 37 25 69 45 101 65 enq 용 Е e 6 06 38 F 70 46 f 102 ack £ 26 66 bel 7 07 39 27 G 71 47 103 67 g bs 8 08 (40 28 н 72 48 h 104 68 ht 9 09) 41 29 Ι 73 49 105 69 10 1f 0A 42 2A J 74 4A i 106 6A vt 11 0B 43 2B K 75 4B k 107 6B ff 12 0C 44 2C L 76 4C 1 108 6C 13 0D 45 2D 77 109 cr Μ 4D 6D 14 78 110 so 0E 46 2E Ν 4E 6E 15 47 79 0F 2F 0 4F 111 6F si 0 dle 16 10 0 48 30 P 80 50 112 70 p 49 31 81 113 dc1 17 11 1 Q 51 q 71 dc2 18 12 2 50 32 82 52 114 72 R r 19 13 51 33 83 53 115 dc3 3 S 73 s 20 14 52 34 т 84 54 t 116 74 dc4 4 nak 21 15 5 53 35 U 85 55 117 75 u 22 54 36 86 56 118 syn 16 6 V v 76 23 17 55 37 87 57 119 etb 7 W W 77 24 18 38 88 120 can 8 56 Х 58 х 78 em 25 19 9 57 39 Y 89 59 У 121 79 sub 26 1A 58 3A Z 90 5A 122 7A : Z esc 27 1B 59 3B 91 5B 123 7B 28 92 fs 1C < 60 3C 5C 124 7C 29 93 125 1D 61 3D] 5D 7D gs = rs 30 1E 62 3E 94 5E 126 7E > 95 us 31 1F ? 63 3F 5F del 127 7F

NOTES: RTL corresponds to execution (after fetch!); JSRR not shown

		$M[BaseR + SEXT(offset6)] \leftarrow SR$		$R7 \leftarrow PC, PC \leftarrow M[ZEXT(trapvect8)]$
E	STR SR, BaseR, offset6	STR 0111 SR BaseR offset6	TRAP trapvect8	TRAP 1111 0000 trapvect8
no		$M[M[PC + SEXT(PCoffset9)]] \leftarrow SR$		R7 ← PC, PC ← PC + SEXT(PCoffset11)
do	STI SR, PCoffset9	STI 1011 SR PCoffset9	JSR PCoffset11	JSR 0100 1 PCoffset11
f E		$M[PC + SEXT(PCoffset9)] \leftarrow SR$		PC ← BaseR
CE	ST SR, PCoffset9	ST 0011 SR PCoffset9	JMP BaseR	JMP 1100 000 BaseR 000000
2		DR ← NOT SR, Setcc		((n AND N) OR (z AND Z) OR (p AND P)): $PC \leftarrow PC + SEXT(PCoffset9)$
20	NOT DR, SR	NOT 1001 DR SR 111111	BR{nzp} PCoffset9	BR 0000 n z p PCoffset9
M		$DR \leftarrow PC + SEXT(PCoffset9)$, Setcc		$DR \leftarrow SR1 \text{ AND SEXT(imm5), Setcc}$
Iid	LEA DR, PCoffset9	LEA 1110 DR PCoffset9	AND DR, SR1, imm5	AND 0101 DR SR1 1 imm5
te		DR ← M[BaseR + SEXT(offset6)], Setcc		DR ← SR1 AND SR2, Setcc
rm	LDR DR, BaseR, offset6	LDR 0110 DR BaseR offset6	AND DR, SR1, SR2	AND 0101 DR SR1 0 00 SR2
\mathbf{E}		DR ← M[M[PC + SEXT(PCoffset9)]], Setcc		DR ← SR1 + SEXT(imm5), Setcc
xaı	LDI DR, PCoffset9	LDI 1010 DR PCoffset9	ADD DR, SR1, imm5	ADD 0001 DR SR1 1 imm5
m		DR ← M[PC + SEXT(PCoffset9)], Setcc		DR ← SR1 + SR2, Setcc
2	LD DR, PCoffset9	LD 0010 DR PCoffset9	ADD DR, SR1, SR2	ADD 0001 DR SR1 0 00 SR2