

# ECE 220: Computer Systems and Programming

## Fall 2020 – Final Exam

1. This is a closed-book, closed-notes exam
2. Absolutely no interaction between students is allowed
3. Illegible handwriting will be graded as incorrect
4. You must put your name and NetID on your submission page
5. **Use a separate page for each question**
6. Submission is only accepted through Gradescope

Question 1 (30 points): \_\_\_\_\_

Question 2 (30 points): \_\_\_\_\_

Question 3 (10 points): \_\_\_\_\_

Question 4 (15 points): \_\_\_\_\_

Question 5 (15 points): \_\_\_\_\_

**Total Score:** \_\_\_\_\_

Your answers should be in the following format. Each question should be on a separate page. **Write out the entire line of code highlighted in yellow.**

On each page, write down the following:

NetID

Q#

(1)

(2)

(3)

...

## Problem 1 (30 points): Linked List

The below program, called “dotProd,” walks through two linked lists to perform a sparse vector dot product. Vectors are simply a 1-dimensional matrix, with only a single row in the case of our lists.

**The vector header is defined as:**

```
typedef struct vector{
    int length;
    node_t *head;
} vect_t;
```

**Nodes in the list are defined as:**

```
typedef struct node{
    double value;
    int col;
    struct node *next;
} node_t;
```

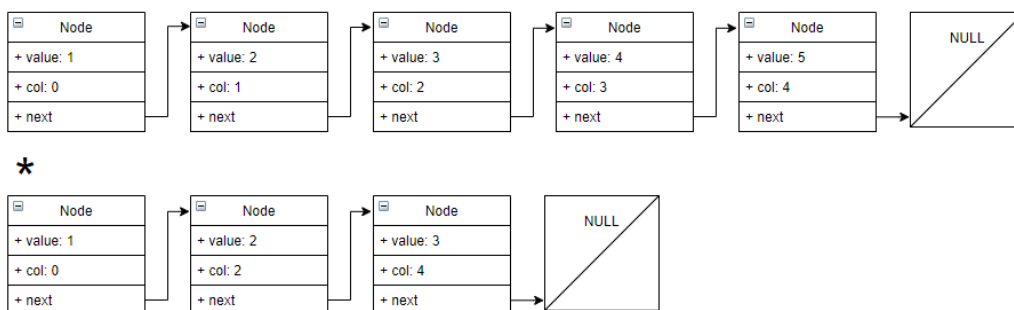
Where col is the column of the given entry. Assume that the list is in order, and that any column without an entry simply has a value of 0, as we assumed for sparse matrices.

The dot product is simply a series of multiply-accumulates. For every column, the values in both vectors are multiplied, and then added into the final result. This means that if a column does not exist in either list, then the product must be zero, and does not change the dot product at all.

**Ex:**  $[1\ 2\ 3\ 4\ 5] * [1\ 0\ 2\ 0\ 3]$

$$\begin{array}{c} [1\ 2\ 3\ 4\ 5] \\ * \\ \begin{array}{c} \Downarrow \Downarrow \Downarrow \Downarrow \Downarrow \\ [1\ 0\ 2\ 0\ 3] \end{array} \\ = 1 + 0 + 6 + 0 + 15 = 22 \end{array}$$

**In LL form:**



## Fill in the blanks to complete the code

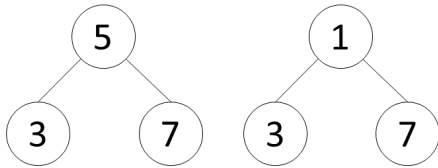
```
1
2
3
4 //sparse dot product
5 //multiply and accumulate both lists
6 //where missing elements are just zeroes
7 double dotProd(vect_t *v1, vect_t *v2){
8
9     //return 0 if dot product cannot be performed
10    //if either vector is nonexistent
11    if( (1) ) return 0;
12    //if the lengths mismatch
13    if( (2) ) return 0;
14
15    //start with no accumulated result
16    double result = 0.0;
17
18    node_t *h1 = v1->head;
19    node_t *h2 = v2->head;
20
21    //while both remaining lists are not NULL
22    while( (3) ){
23
24        //if they match in col
25        if( (4) ){
26            //multiply and accumulate
27            (5) += (6) ;
28            h1 = h1->next;
29            h2 = h2->next;
30        }
31        else if( (7) ){ //if h1 precedes h2
32            h1 = (8) ; //move h1 forward
33        }
34        else{
35            h2 = (9) ; //move h2 forward
36        }
37    }
38    return result;
39 }
40 }
```

## Problem 2 (30 points): Binary Search Tree

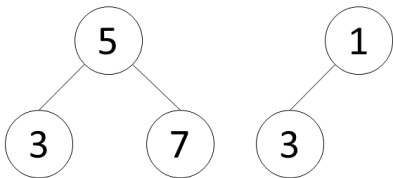
### Part A.

The program, called “treematch,” uses a post-order traversal of a tree to compare two trees and determine whether

- the trees are exactly the same, including their values
- the trees’ structures are the same, but the values differ
- the trees’ structures differ



Above is an example of two trees with different values, but the same structure. The root node for the left tree is 5, while the root node for the right tree is 1.



For this example, the structures differ, so the fact that the values differ does not matter. The root node of the left tree has two children, but the root node of the right tree has only a left child.

**The returned value of `treematch()` is an enum, defined as:**

```
typedef enum{
    SAME_ALL = 0
    SAME_STRUCT_DIFF_VALUE = 1
    DIFF_STRUCT = 2
}match_t;
```

`SAME_ALL` means that both the structure and the values are the same.

`SAME_STRUCT_DIFF_VALUE` means the structure is the same, but the values differ.

`DIFF_STRUCT` means that the structures differ (the values don’t matter in this case).

**The tree node is defined as:**

```
typedef struct node{
    int value;
    struct node *parent, *left, *right;
}node_t;
```

**You may assume:**

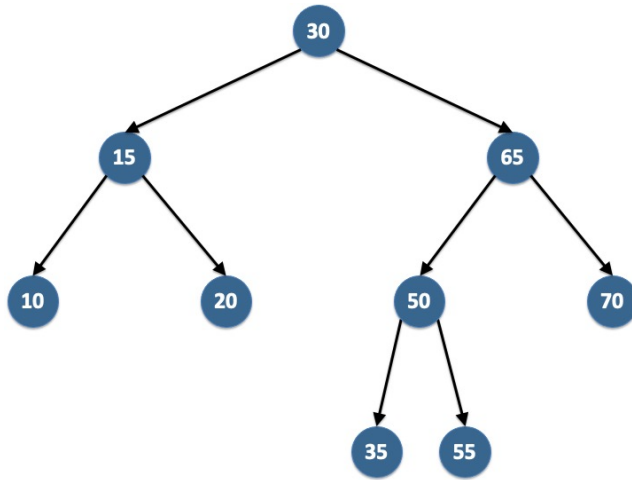
1. That the tree is formed correctly (e.g. that parent and child links are correct).

## Fill in the blanks to complete the code

```
1 #include <stdio.h>
2 #include <stdlib.h>
3 #include "treematch.h"
4
5 match_t treematch(node_t *roota, node_t *rootb) {
6     //if both null
7     if ( (1) _____ ) {
8         //consider as having same value and structure
9         return SAME_ALL;
10    }
11    //check if structure matches
12    if ( (2) _____ ){
13        //recursively check subtrees
14        match_t mleft = treematch( (3) _____ );
15        match_t mright = treematch( (4) _____ );
16
17        //if either subtree has a different structure
18        if ( (5) _____ ) {
19            return DIFF_STRUCT ;
20        }
21        //if either subtree has a different value
22        else if ( (6) _____ ) {
23            return SAME_STRUCT_DIFF_VALUE ;
24        }
25        //remaining: the subtrees are the same
26        //check the values in the current nodes
27        if ( (7) _____ ) {
28            return SAME_ALL;
29        }
30        else {
31            return (8) _____ ;
32        }
33    }
34    //remaining case: one pointer exists, the other doesn't
35    return (9) _____ ;
36 }
```

**Part B.**

Given the binary search tree below, write out the **sequence of nodes visited during a pre-order traversal**.



**Your Answer:**

**Part C.**

Where should a new node with the value '62' be inserted into the binary search tree shown in Part B? **Draw the entire tree for your answer.**

**Your Answer:**

### Problem 3 (10 points): C++

In this problem, we are going to implement some geometric computation with the concept of OOP.

Point is a class contains 2 data members in public: x and y in double, that is the 2D coordinate.

There is a base class Polygon that contains a member function `get_area`, which is going to compute the area for the polygon.

The class Triangle, public derived from the class Polygon, that contains 3 data members (3 Point): `point1`, `point2`, `point3` in private. Please complete the class Triangle for its constructor, copy constructor, access function, and the `get_area` function.

**List of the functions you need to complete:**

**Default constructor:** initialize `point1` as (0, 0), `point2` as (1, 0), `point3` as (0,1).

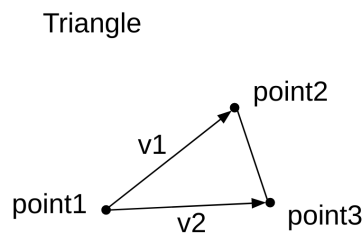
**Copy constructor:** copy the points value from the given triangle.

**Constructor:** initialize the points' value with the given points, `tp1`, `tp2`, `tp3`.

**Access function:** let the public domain be able to access the private data: `point1`, `point2`, `point3`.

**get\_area function:** compute the area of the triangle by computing the cross product of two vectors.

- Assign `v1` as a vector from `p1` to `p2` by subtracting `p1`'s coordinates from `p2`; `v2` as a vector from `p1` to `p3` by subtracting `p1`'s coordinates from `p3`.
- Compute the cross product of `v1` and `v2`. (Hint: cross product:  $(x_1, y_1) \times (x_2, y_2) = x_1 y_2 - x_2 y_1$ )
- Get the absolute of this product and divide it by 2.



```

#include <iostream>
using namespace std;
class Point {
private:
    double x, y;
public:
    void set_val(const double & tx, const double & ty){
        x = tx;
        y = ty;
    }
    Point():x(0),y(0){};
    Point(const Point & p):x(p.x()), y(p.y()){};
    Point(const double tx, const double ty):x(tx), y(ty){};
    ~Point(){};
    const double x() const{ return x; }
    const double y() const{ return y; }
};

class Polygon{
private:
public:
    virtual double get_area() const = 0 ;
};

class Triangle: public Polygon {
private:
    Point point1, point2, point3;
public:
    const Point p1() const;
    const Point p2() const;
    const Point p3() const;
    Triangle();
    Triangle(const Triangle &t);
    Triangle(const Point & p1, const Point & p2, const Point & p3 );
    ~Triangle(){};
    double get_area() const;
};

Triangle::Triangle(){ // default ctor
// initial p1 = (0,0), p2 = (1, 0), p3 = (0,1)
(1) _____
(2) _____
(3) _____
}

Triangle::Triangle(const Triangle & t){ // copy ctor
(4) _____
(5) _____
(6) _____
}

```



```

}
Triangle::Triangle(const Point & p1, const Point & p2, const Point & p3 ){ // ctor w/ given initial value
(7) _____
(8) _____
(9) _____
}
//Access function
const Point Triangle::p1() const{
return (10) _____
}
const Point Triangle::p2() const{
return (11) _____
}
const Point Triangle::p3() const{
return (12) _____
}
// compute area function
double Triangle::get_area() const{
Point v1( (13) _____ ); // vector p1 --> p2
Point v2( (14) _____ ); // vector p1 --> p3
double cross_product = (15) _____
if (cross_product < 0) { // compute the absolute value
cross_product = -cross_product;
}
return 0.5*cross_product;
}

int main(){
Polygon* ptr;

Triangle tri0;
ptr = &tri0;
std::cout << "tri 0 area : "<<ptr->get_area()<<std::endl;

Point p1(1,1);
Point p2(4,4);
Point p3(4,1);
Triangle tri1(p1, p2, p3);
ptr = &tri1;
std::cout << "tri1 area : "<<ptr->get_area()<<std::endl;

Triangle tri2(tri1);
ptr = &tri2;
std::cout << "tri2 area : "<<ptr->get_area()<<std::endl;

return 0;
}

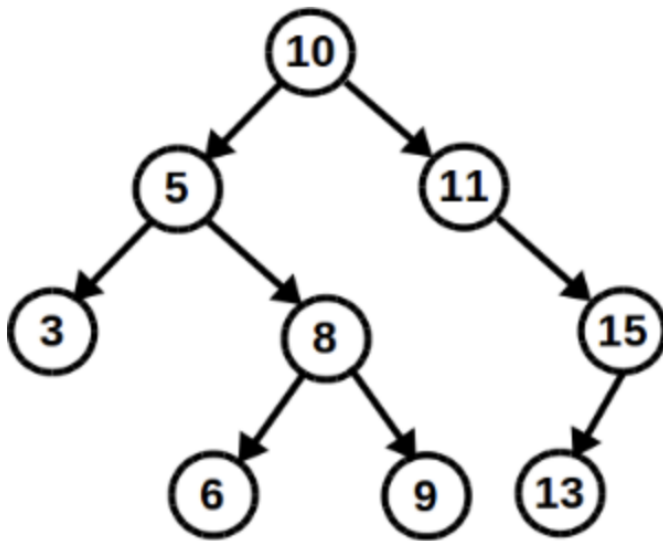
```

## Problem 4 (15 points): C to LC-3 Conversion

In this problem, you will translate a recursive C function **SumPostOrder**, which traverses the binary tree in post-order and adds the value of each node in the tree and prints the current sum, to LC3.

You can assume each memory location stores 16 bits. Both integers and pointers are also 16 bits.

**Example:**



**Input:** root of binary tree    **Output:** 80

**C code:**

```
typedef struct Node{
    int value;
    Node* left;
    Node* right;
}

int SumPostOrder(Node* node) {
    int sum = 0;
    /* Base case */
    if (node == NULL) {
        return sum;
    }

    /* Recursive calls */
    sum += SumPostOrder((node->left));
    sum += SumPostOrder(node->right);
    sum += node->value;
    return sum;
}
```

Recalled that a function's activation record has the following format:

Local Variables
Caller's Frame Pointer
Return Address
Return Value
Arguments

**Register map for this problem:**

<b>R0</b>	NODE
<b>R1</b>	SUM
<b>R2</b>	TEMP
<b>R5</b>	FRAME POINTER
<b>R6</b>	STACK POINTER
<b>R7</b>	RETURN ADDRESS

**Requirements:**

You **MUST** follow the run-time stack convention presented in lectures.

You **MUST** finish your code within the number of provided lines.

You **MUST** use the register map as specified above. No other registers may be used.

```
SUM_POST_ORDER
; callee set-up
(1)                     
(2)                     
(3)                     
(4)                     
;
; initialize sum
(5)                     
;
; check base case
(6)                      ; load NODE
(7)                      ; check NODE==NULL
;
; SumPostOrder(node->left)
(8)                      ; load NODE->left to TEMP
(9)                      ; update stack pointer
(10)                    ; push argument
JSR SUM_POST_ORDER
(11)                      ; store return value in TEMP
(12)                      ; add TEMP to SUM
(13)                      ; stack teardown
;
; SumPostOrder(node->right)
; second recursive call omitted for simplicity
;
; add current node's value to sum
(14)                      ; load NODE->value to TEMP
(15)                      ; add NODE->value to SUM
;
DONE
(16)                      ; store return value
; callee tear-down
; omitted here for simplicity
RET
```

## Problem 4 (15 points): Concepts

### Part A & B

Recall in MP8 - 2048, we worked with a dynamically allocated game struct:

```
typedef int cell;
typedef struct{
    int rows;
    int cols;
    cell* cells;
    int score;
} game;
game* make_game(int rows, int cols){
    game* mygame = (game*)malloc(sizeof(game));
    mygame->cells = malloc(rows * cols * sizeof(cell));
    // rest of the function is omitted
}
```

Part A: Why is it necessary that the following function is executed after each game finishes?

```
void destroy_game(game* cur_game){
    free(cur_game->cells);
    free(cur_game);
    return NULL;
}
```

**Your Answer (use no more than 20 words):**

Part B: It is not possible to adjust the size of the **cells** array after it has been created. (True or False)

**Your Answer:**

True

False

## Part C

Given the C++ classes below, answer the question in Part C. Note that the code is different from what was provided for MP12.

```
1 class Number{
2     public:
3         double magnitude, phase;
4         Number();
5         // more functions omitted
6         virtual Number operator + (const Number& arg) = 0;
7 };
8 class RealNumber : private Number{
9     public:
10        double real_component;
11        RealNumber();
12        RealNumber operator + (const RealNumber& arg);
13        // more functions omitted
14 };
```

In line 12, how is the argument **arg** passed to the operator overload function? Choose the correct answer from the following options.

1. By value

2. By pointer

3. By reference

**Your Answer for Part C:**

**Table E.2 The Standard ASCII Table**

ASCII			ASCII			ASCII			ASCII		
Character	Dec	Hex	Character	Dec	Hex	Character	Dec	Hex	Character	Dec	Hex
nul	0	00	sp	32	20	@	64	40	`	96	60
soh	1	01	!	33	21	A	65	41	a	97	61
stx	2	02	"	34	22	B	66	42	b	98	62
etx	3	03	#	35	23	C	67	43	c	99	63
eot	4	04	\$	36	24	D	68	44	d	100	64
enq	5	05	%	37	25	E	69	45	e	101	65
ack	6	06	&	38	26	F	70	46	f	102	66
bel	7	07	'	39	27	G	71	47	g	103	67
bs	8	08	(	40	28	H	72	48	h	104	68
ht	9	09	)	41	29	I	73	49	i	105	69
lf	10	0A	*	42	2A	J	74	4A	j	106	6A
vt	11	0B	+	43	2B	K	75	4B	k	107	6B
ff	12	0C	,	44	2C	L	76	4C	l	108	6C
cr	13	0D	-	45	2D	M	77	4D	m	109	6D
so	14	0E	.	46	2E	N	78	4E	n	110	6E
si	15	0F	/	47	2F	O	79	4F	o	111	6F
dle	16	10	0	48	30	P	80	50	p	112	70
dc1	17	11	1	49	31	Q	81	51	q	113	71
dc2	18	12	2	50	32	R	82	52	r	114	72
dc3	19	13	3	51	33	S	83	53	s	115	73
dc4	20	14	4	52	34	T	84	54	t	116	74
nak	21	15	5	53	35	U	85	55	u	117	75
syn	22	16	6	54	36	V	86	56	v	118	76
etb	23	17	7	55	37	W	87	57	w	119	77
can	24	18	8	56	38	X	88	58	x	120	78
em	25	19	9	57	39	Y	89	59	y	121	79
sub	26	1A	:	58	3A	Z	90	5A	z	122	7A
esc	27	1B	;	59	3B	[	91	5B	{	123	7B
fs	28	1C	<	60	3C	\	92	5C		124	7C
gs	29	1D	=	61	3D	]	93	5D	}	125	7D
rs	30	1E	>	62	3E	^	94	5E	~	126	7E
us	31	1F	?	63	3F	_	95	5F	del	127	7F

NOTES: RTL corresponds to execution (after fetch!); JSRR not shown

LC-3 Instructions

ADD	<table border="1"><tr><td>0001</td><td>DR</td><td>SR1</td><td>0</td><td>00</td><td>SR2</td></tr></table>	0001	DR	SR1	0	00	SR2	ADD DR, SR1, SR2	LD	<table border="1"><tr><td>0010</td><td>DR</td><td colspan="4">PCOffset9</td></tr></table>	0010	DR	PCOffset9				LD DR, PCOffset9
0001	DR	SR1	0	00	SR2												
0010	DR	PCOffset9															
		$DR \leftarrow SR1 + SR2, Setcc$			$DR \leftarrow M[PC + SEXT(PCOffset9)], Setcc$												
ADD	<table border="1"><tr><td>0001</td><td>DR</td><td>SR1</td><td>1</td><td colspan="2">imm5</td></tr></table>	0001	DR	SR1	1	imm5		ADD DR, SR1, imm5	LDI	<table border="1"><tr><td>1010</td><td>DR</td><td colspan="4">PCOffset9</td></tr></table>	1010	DR	PCOffset9				LDI DR, PCOffset9
0001	DR	SR1	1	imm5													
1010	DR	PCOffset9															
		$DR \leftarrow SR1 + SEXT(imm5), Setcc$			$DR \leftarrow M[M[PC + SEXT(PCOffset9)]], Setcc$												
AND	<table border="1"><tr><td>0101</td><td>DR</td><td>SR1</td><td>0</td><td>00</td><td>SR2</td></tr></table>	0101	DR	SR1	0	00	SR2	AND DR, SR1, SR2	LDR	<table border="1"><tr><td>0110</td><td>DR</td><td>BaseR</td><td colspan="3">offset6</td></tr></table>	0110	DR	BaseR	offset6			LDR DR, BaseR, offset6
0101	DR	SR1	0	00	SR2												
0110	DR	BaseR	offset6														
		$DR \leftarrow SR1 \text{ AND } SR2, Setcc$			$DR \leftarrow M[BaseR + SEXT(offset6)], Setcc$												
AND	<table border="1"><tr><td>0101</td><td>DR</td><td>SR1</td><td>1</td><td colspan="2">imm5</td></tr></table>	0101	DR	SR1	1	imm5		AND DR, SR1, imm5	LEA	<table border="1"><tr><td>1110</td><td>DR</td><td colspan="4">PCOffset9</td></tr></table>	1110	DR	PCOffset9				LEA DR, PCOffset9
0101	DR	SR1	1	imm5													
1110	DR	PCOffset9															
		$DR \leftarrow SR1 \text{ AND } SEXT(imm5), Setcc$			$DR \leftarrow PC + SEXT(PCOffset9), Setcc$												
BR	<table border="1"><tr><td>0000</td><td>n</td><td>z</td><td>p</td><td colspan="2">PCOffset9</td></tr></table>	0000	n	z	p	PCOffset9		BR{nzp} PCOffset9	NOT	<table border="1"><tr><td>1001</td><td>DR</td><td>SR</td><td colspan="3">11111</td></tr></table>	1001	DR	SR	11111			NOT DR, SR
0000	n	z	p	PCOffset9													
1001	DR	SR	11111														
		$((n \text{ AND } N) \text{ OR } (z \text{ AND } Z) \text{ OR } (p \text{ AND } P)):$ $PC \leftarrow PC + SEXT(PCOffset9)$			$DR \leftarrow \text{NOT } SR, Setcc$												
JMP	<table border="1"><tr><td>1100</td><td>000</td><td>BaseR</td><td colspan="3">000000</td></tr></table>	1100	000	BaseR	000000			JMP BaseR	ST	<table border="1"><tr><td>0011</td><td>SR</td><td colspan="4">PCOffset9</td></tr></table>	0011	SR	PCOffset9				ST SR, PCOffset9
1100	000	BaseR	000000														
0011	SR	PCOffset9															
		$PC \leftarrow BaseR$			$M[PC + SEXT(PCOffset9)] \leftarrow SR$												
JSR	<table border="1"><tr><td>0100</td><td>1</td><td colspan="4">PCOffset11</td></tr></table>	0100	1	PCOffset11				JSR PCOffset11	STI	<table border="1"><tr><td>1011</td><td>SR</td><td colspan="4">PCOffset9</td></tr></table>	1011	SR	PCOffset9				STI SR, PCOffset9
0100	1	PCOffset11															
1011	SR	PCOffset9															
		$R7 \leftarrow PC, PC \leftarrow PC + SEXT(PCOffset11)$			$M[M[PC + SEXT(PCOffset9)]] \leftarrow SR$												
TRAP	<table border="1"><tr><td>1111</td><td>0000</td><td colspan="4">trapvect8</td></tr></table>	1111	0000	trapvect8				TRAP trapvect8	STR	<table border="1"><tr><td>0111</td><td>SR</td><td>BaseR</td><td colspan="3">offset6</td></tr></table>	0111	SR	BaseR	offset6			STR SR, BaseR, offset6
1111	0000	trapvect8															
0111	SR	BaseR	offset6														
		$R7 \leftarrow PC, PC \leftarrow M[ZEXT(trapvect8)]$			$M[BaseR + SEXT(offset6)] \leftarrow SR$												

End of ECE 220 Final Exam