

ECE 220: Computer Systems and Programming

Spring 2021 – Final Exam

May 10, 2021

1. This is a closed-book, closed-notes exam
2. Absolutely no interaction between students is allowed
3. Illegible handwriting will be graded as incorrect
4. You must put your name and NetID on your submission page
5. **Use a separate page for each question**
6. Submission is only accepted through Gradescope
7. **Your CBTF time stamp is checked against your Gradescope submission time stamp. Any unusual activity will be reported to the college for infraction of academic integrity.**

Question 1 (24 points): _____

Question 2 (30 points): A)____; B)____; C)____

Question 3 (16 points): _____

Question 4 (12 points): _____

Question 5 (18 points): A)____; B)____; C)____; D)____

Total Score: _____

Write down your answers in the following format. Each question should be on a separate page. **Tag each question on your Gradescope submission.**

Name: _____ NetID: _____

Q1 (**write down the # and entire line of code highlighted in yellow**)

Q2 (**write down the # and entire line of code highlighted in yellow**)

Part A; Part B; Part C

Q3 (**write down the # and entire line of code highlighted in yellow**)

Q4 (**write down the # and entire line of code highlighted in yellow**)

Q5

Part A; Part B; Part C; Part D

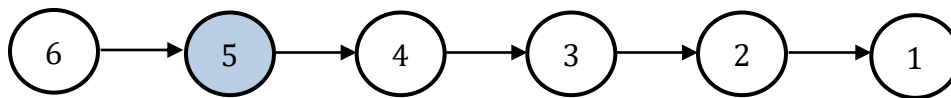
Problem 1 (24 points): Linked List

Given the head of a linked list, and an integer k . Write a C function that removes the k th node (the list is 1-indexed) from the head and insert the removed node as the $(2k)$ -th node of the linked list. Assume $k > 1$ and the length of the linked list is larger than k . If the length of the linked list is smaller than $2k$, please insert the removed node to the end of the linked list.

Example

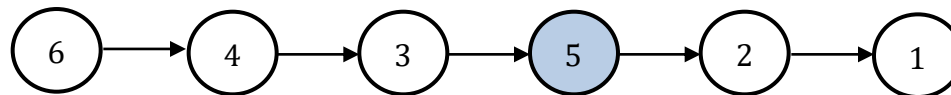
Input:

[6,5,4,3,2,1], $k=2$



Output:

[6,4,3,5,2,1]



```

#include <stddef.h>
#include <stdlib.h>
#include <stdio.h>

typedef struct ListNode node;

struct ListNode {
    int data;
    struct ListNode *next;
};

void KthNode(node* head, int k){
    node *cur = head;
    int i;
    /* traverse the list to find the (k-1)-th node */
    for(i=1; i < k - 1; i++){
        (1) _____;
    }

    /* save the pointer that points to the (k)-th node */
    node *k_ptr = (2) _____;

    /* update pointer to "remove" the (k)-th node */
    cur->next = (3) _____;

    /* continue to traverse the list to find the right place for
       inserting the (k)-th node */
    i = 0;
    while(cur->next != NULL){
        cur = (4) _____;
        i = (5) _____;
        if(i == k)
            (6) _____;
    }

    /* insert the (k)-th node after 'cur' node */
    node *tmp = cur->next;
    cur->next = (7) _____;
    k_ptr->next = (8) _____;
}

```

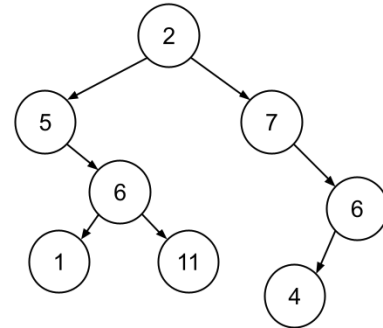
Problem 2 (30 points): Binary Tree

This problem has Parts A, B, and C.

Part A: Given a binary tree, write a recursive C function to find its largest node and return a pointer to it. Please note that the tree does not have duplicate nodes.

Example:

For the given binary tree, the function should return a pointer to the node contains the data '11'.

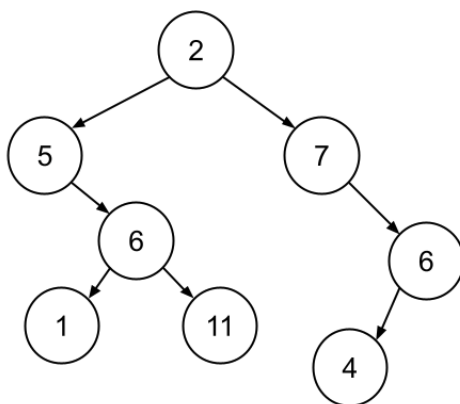


Part B: Write a C function to insert a new node as the left child of the largest node (found in Part A). If there already exists a left child, insert the new node as the right child. If the left and right child already exists, insert the new node as a left child anyway. Move the existing left subtree of the largest node to be the left child of the new node that is inserted.

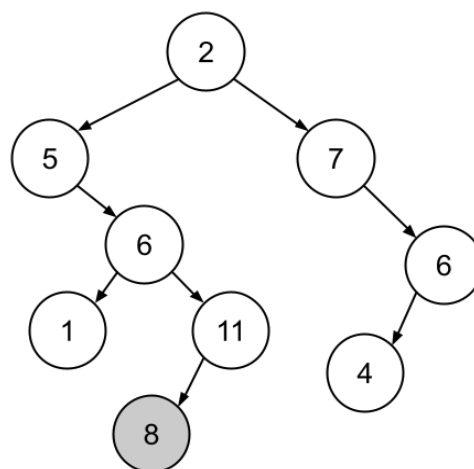
Examples:

In all the examples below, the maximum value found from part A is 11 and the new node to insert is 8.

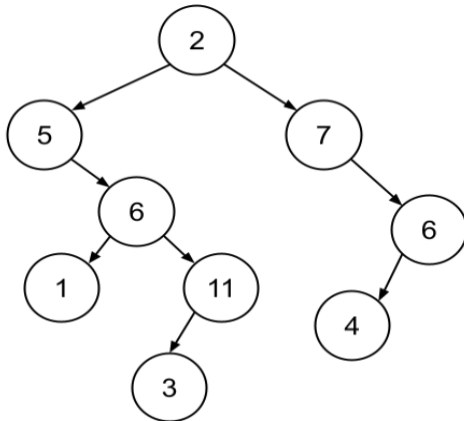
1. Initial Tree:



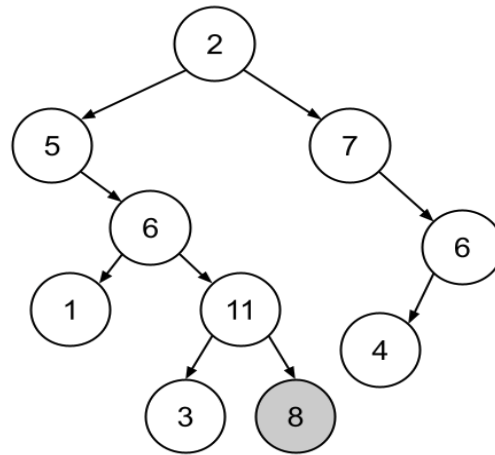
After Insertion:



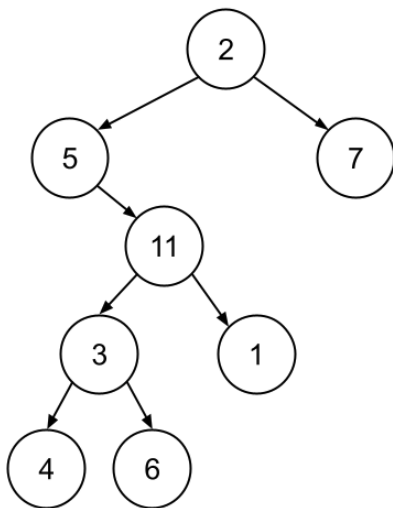
2. Initial Tree:



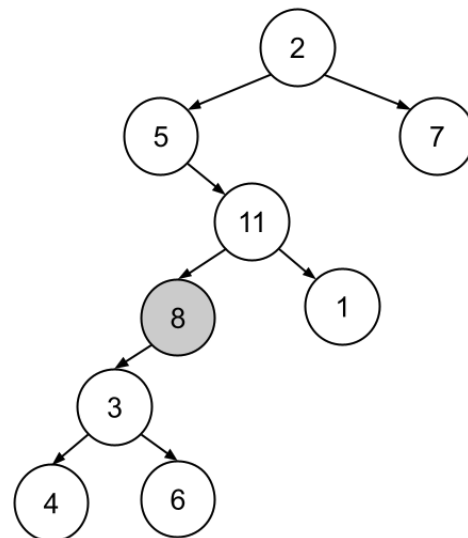
After Insertion:



3. Initial Tree:



After Insertion:



For Part A and Part B, fill in the blanks below.

```

/* tree node */
struct Node {
    int data;
    struct Node *left, *right;
};
  
```

```

/* Part A: find the maximum node */
struct Node* findMax(struct Node* root){
    /* Base case */
    if ((1) )
        return NULL;

    int lres=0;
    int rres=0;

    /* Return maximum of 3 values: 1) Root's data, 2) Max
    in Left Subtree, 3) Max in right subtree */

    int res = root->data;

    /* Recursion for the left subtree */
    struct Node* left_max = (2) ;
    if (left_max!=NULL)
        lres=left_max->data;

    /* Recursion for the right subtree */
    struct Node* right_max = (3) ;
    if (right_max!=NULL)
        rres=right_max->data;

    /* Compare lres, rres and res then return the ptr to
    the max node */
    /* if left subtree max node's data is greater than
    that of right subtree's max and root */
    if ((4) ){
        return left_max;
    }
    /* if right subtree max node's data greater than that
    of root */
    if ((5) ){
        return right_max;
    }
    return root;
}

```

```

/* Part B: Inserts a new node */
/* Assume helper function 'struct Node* newnode(int val);' is
given and can be used to dynamically allocate a new node, set its
data to val, and left and right pointers to NULL */
void insertNode(struct Node* max_node, int val){
    if ((6) ) {
        printf("Cannot Insert New Node!");
        return;
    }

    /* Case 1: max_node doesn't have a left child */
    if ((7) ) {
        (8) ;
        printf("Node inserted as left child.\n");
        return;
    }

    /* Case 2: max_node has a left child but no right child */
    if ((9) ) {
        (10) ;
        printf("Node inserted as right child.\n");
        return;
    }

    /* Case 3: max_node has both left and right children */
    struct Node* temp = max_node->left;
    max_node->left = (11) ;
    (12) ;
    printf("Node inserted as left child. Inserted node now has a
left subtree.\n");
}

/* Main Function */
int main(){
    /* creation of the tree is omitted*/

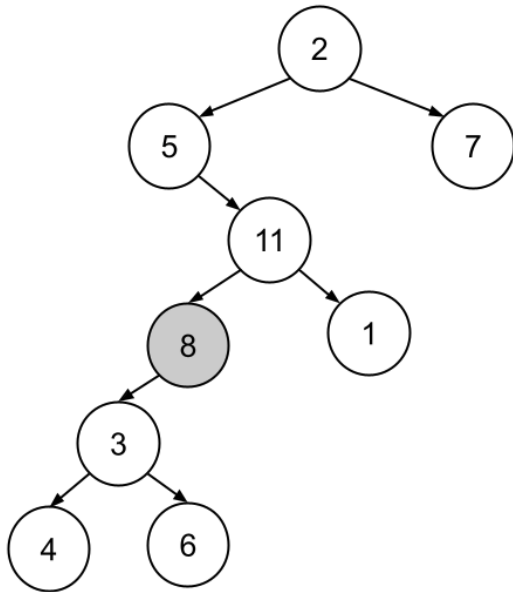
    /* Function call for part A. Assume root to be the root of
the binary tree */
    struct Node* max_node=findMax(root);

    if (root!=NULL)
        printf("Maximum node is %d \n", max_node->data);

    /* Function call for part B. Assume 8 is to be inserted */
    insertNode(max_node, 8);
    return 0;
}

```

Part C: Suppose the binary tree after insertion looks like this:



What sequence would an in-order traversal (left, root, right) produce?

Your Answer: _____

Problem 3 (16 points): C to LC-3 Conversion

In the following, **alternateList(struct Node* head)** represents a C function that prints alternate nodes of the given Linked List, first from head to tail, and then from tail to head. If Linked List has an even number of nodes, then **alternateList()** skips the last node.

Examples:

If given a linked list: 1->2->3->4->5, it prints 1 3 5 5 3 1.

If given a linked list: 1->2->3->4->5->6, it prints 1 3 5 5 3 1.

Replicate the given C function **alternateList(struct Node* head)** using LC3. Code snippet is provided. You just need to fill in the blanks.

```
_____ C Code _____

/* A Linked List Node */
struct Node{
    int data;
    struct Node *next;
};

/* Recursive function - this function has no local variables */
void alternateList(struct Node* head){
    if(head == NULL)
        return;
    printf("%d ", head->data);

    if(head->next != NULL )
        alternateList(head->next->next);
    printf("%d ", head->data);
}
```

Recall that a function's activation record has the following format:

Local Variables
Caller's Frame Pointer
Return Address
Return Value
Arguments

You **MUST** use and conform to the **LC-3 calling conventions** we have described in class. You may assume each data type will only occupy one memory location in LC-3.

```
; R6 is the stack pointer. R5 is the frame pointer.  
; R7 contains the return address.
```

```
ALTERNATE_LIST
```

```
; complete callee set-up (push bookkeeping information and  
; local variables) and comment on each line  
(1) _____; reserve spaces (update stack pointer)  
(2) _____; store return address  
(3) _____; store caller's frame pointer  
(4) _____; update frame pointer
```

```
; Load head to R0
```

```
(5) _____  
; Check if head is NULL  
(6) _____  
(7) BR _____ FINISH
```

```
; printing omitted
```

```
; Check if head->next is NULL
```

```
(8) _____; load head->next to R2  
(9) BR _____ PRINT
```

```
RECURSIVE
```

```
(10) _____; load head->next->next to R4  
; complete caller set-up (push arguments)  
(11) _____; store argument to run-time stack  
(12) _____; update stack pointer  
; make recursive call  
(13) _____  
; caller tear-down omitted
```

```
PRINT
```

```
; printing omitted
```

```
FINISH
```

```
; callee teardown
```

```
(14) _____; restore caller's frame pointer  
(15) _____; restore return address  
(16) _____; update stack pointer  
RET
```

Problem 4 (12 points): C++

The following simple C++ program declares a class, creates some objects, prints them, and performs a “+” operation. Fill out the blanks in main.cpp for the program to match the provided output.

The expected output should look like the following:

```
This house has 5 rooms.  
This house has 10 rooms.  
This house has 15 rooms.
```

main.cpp

```
#include<iostream>
using namespace std;
class house {
    int roomCount;
public :
    house() {}
    house(int r){

        _____
    }

    int getRoomCount(void){

        _____
    }

    void setRoomCount(int newRoomCount){

        _____
    }

    void houseInfo() {
        cout << "This house has "<<roomCount<< " rooms."<<endl;
    }

    house operator + (const house & other){

        _____
        _____
        _____

    }
};
```

```
int main() {  
    house a(5);  
    house b(10);  
    a.houseInfo();  
    b.houseInfo();  
    house c = a+b;  
    c.houseInfo();  
    return 0;  
}
```

Problem 5 (18 points): Concepts

Part A

Consider the following code:

```
char *name;  
name = (char *)malloc(32*sizeof(char));  
realloc(name, 16*sizeof(char));  
free(name);
```

Assume memory allocations are successful and each char is 8-bit long, how many bytes of memory is deallocated when we call free?

Your Answer: _____ bytes

Part B

What would cause a memory leak?

- (a) Run out of space in the run-time stack
- (b) Run out of space in the heap
- (c) Lost a pointer to a block of dynamically allocated memory
- (d) Access memory location that is not allowed

Your Answer: **A** **B** **C** **D**

Part C

Determine whether the following statement is true or false:

A stack can be implemented by using an array or a linked list.

Your Answer: **True** **False**

Part D

Determine whether the following statement is true or false:

In C++, a virtual function table in each class stores its virtual functions.

Your Answer: **True** **False**

Table E.2 The Standard ASCII Table

ASCII			ASCII			ASCII			ASCII		
Character	Dec	Hex	Character	Dec	Hex	Character	Dec	Hex	Character	Dec	Hex
nul	0	00	sp	32	20	@	64	40	`	96	60
soh	1	01	!	33	21	A	65	41	a	97	61
stx	2	02	"	34	22	B	66	42	b	98	62
etx	3	03	#	35	23	C	67	43	c	99	63
eot	4	04	\$	36	24	D	68	44	d	100	64
enq	5	05	%	37	25	E	69	45	e	101	65
ack	6	06	&	38	26	F	70	46	f	102	66
bel	7	07	'	39	27	G	71	47	g	103	67
bs	8	08	(40	28	H	72	48	h	104	68
ht	9	09)	41	29	I	73	49	i	105	69
lf	10	0A	*	42	2A	J	74	4A	j	106	6A
vt	11	0B	+	43	2B	K	75	4B	k	107	6B
ff	12	0C	,	44	2C	L	76	4C	l	108	6C
cr	13	0D	-	45	2D	M	77	4D	m	109	6D
so	14	0E	.	46	2E	N	78	4E	n	110	6E
si	15	0F	/	47	2F	O	79	4F	o	111	6F
dle	16	10	0	48	30	P	80	50	p	112	70
dc1	17	11	1	49	31	Q	81	51	q	113	71
dc2	18	12	2	50	32	R	82	52	r	114	72
dc3	19	13	3	51	33	S	83	53	s	115	73
dc4	20	14	4	52	34	T	84	54	t	116	74
nak	21	15	5	53	35	U	85	55	u	117	75
syn	22	16	6	54	36	V	86	56	v	118	76
etb	23	17	7	55	37	W	87	57	w	119	77
can	24	18	8	56	38	X	88	58	x	120	78
em	25	19	9	57	39	Y	89	59	y	121	79
sub	26	1A	:	58	3A	Z	90	5A	z	122	7A
esc	27	1B	;	59	3B	[91	5B	{	123	7B
fs	28	1C	<	60	3C	\	92	5C		124	7C
gs	29	1D	=	61	3D]	93	5D	}	125	7D
rs	30	1E	>	62	3E	^	94	5E	_	126	7E
us	31	1F	?	63	3F	_	95	5F	del	127	7F

NOTES: RTL corresponds to execution (after fetch!); JSRR not shown

LC-3 Instructions

ADD	<table border="1"><tr><td>0001</td><td>DR</td><td>SR1</td><td>0</td><td>00</td><td>SR2</td></tr></table>	0001	DR	SR1	0	00	SR2	ADD DR, SR1, SR2	LD	<table border="1"><tr><td>0010</td><td>DR</td><td colspan="4">PCOffset9</td></tr></table>	0010	DR	PCOffset9				LD DR, PCOffset9
0001	DR	SR1	0	00	SR2												
0010	DR	PCOffset9															
		$DR \leftarrow SR1 + SR2, Setcc$			$DR \leftarrow M[PC + SEXT(PCOffset9)], Setcc$												
ADD	<table border="1"><tr><td>0001</td><td>DR</td><td>SR1</td><td>1</td><td colspan="2">imm5</td></tr></table>	0001	DR	SR1	1	imm5		ADD DR, SR1, imm5	LDI	<table border="1"><tr><td>1010</td><td>DR</td><td colspan="4">PCOffset9</td></tr></table>	1010	DR	PCOffset9				LDI DR, PCOffset9
0001	DR	SR1	1	imm5													
1010	DR	PCOffset9															
		$DR \leftarrow SR1 + SEXT(imm5), Setcc$			$DR \leftarrow M[M[PC + SEXT(PCOffset9)]], Setcc$												
AND	<table border="1"><tr><td>0101</td><td>DR</td><td>SR1</td><td>0</td><td>00</td><td>SR2</td></tr></table>	0101	DR	SR1	0	00	SR2	AND DR, SR1, SR2	LDR	<table border="1"><tr><td>0110</td><td>DR</td><td>BaseR</td><td colspan="3">offset6</td></tr></table>	0110	DR	BaseR	offset6			LDR DR, BaseR, offset6
0101	DR	SR1	0	00	SR2												
0110	DR	BaseR	offset6														
		$DR \leftarrow SR1 \text{ AND } SR2, Setcc$			$DR \leftarrow M[BaseR + SEXT(offset6)], Setcc$												
AND	<table border="1"><tr><td>0101</td><td>DR</td><td>SR1</td><td>1</td><td colspan="2">imm5</td></tr></table>	0101	DR	SR1	1	imm5		AND DR, SR1, imm5	LEA	<table border="1"><tr><td>1110</td><td>DR</td><td colspan="4">PCOffset9</td></tr></table>	1110	DR	PCOffset9				LEA DR, PCOffset9
0101	DR	SR1	1	imm5													
1110	DR	PCOffset9															
		$DR \leftarrow SR1 \text{ AND } SEXT(imm5), Setcc$			$DR \leftarrow PC + SEXT(PCOffset9), Setcc$												
BR	<table border="1"><tr><td>0000</td><td>n</td><td>z</td><td>p</td><td colspan="2">PCOffset9</td></tr></table>	0000	n	z	p	PCOffset9		BR(nzp) PCOffset9	NOT	<table border="1"><tr><td>1001</td><td>DR</td><td>SR</td><td colspan="3">11111</td></tr></table>	1001	DR	SR	11111			NOT DR, SR
0000	n	z	p	PCOffset9													
1001	DR	SR	11111														
		$((n \text{ AND } N) \text{ OR } (z \text{ AND } Z) \text{ OR } (p \text{ AND } P)):$ $PC \leftarrow PC + SEXT(PCOffset9)$			$DR \leftarrow \text{NOT } SR, Setcc$												
JMP	<table border="1"><tr><td>1100</td><td>000</td><td>BaseR</td><td colspan="3">000000</td></tr></table>	1100	000	BaseR	000000			JMP BaseR	ST	<table border="1"><tr><td>0011</td><td>SR</td><td colspan="4">PCOffset9</td></tr></table>	0011	SR	PCOffset9				ST SR, PCOffset9
1100	000	BaseR	000000														
0011	SR	PCOffset9															
		$PC \leftarrow BaseR$			$M[PC + SEXT(PCOffset9)] \leftarrow SR$												
JSR	<table border="1"><tr><td>0100</td><td>1</td><td colspan="4">PCOffset11</td></tr></table>	0100	1	PCOffset11				JSR PCOffset11	STI	<table border="1"><tr><td>1011</td><td>SR</td><td colspan="4">PCOffset9</td></tr></table>	1011	SR	PCOffset9				STI SR, PCOffset9
0100	1	PCOffset11															
1011	SR	PCOffset9															
		$R7 \leftarrow PC, PC \leftarrow PC + SEXT(PCOffset11)$			$M[M[PC + SEXT(PCOffset9)]] \leftarrow SR$												
TRAP	<table border="1"><tr><td>1111</td><td>0000</td><td colspan="4">trapvect8</td></tr></table>	1111	0000	trapvect8				TRAP trapvect8	STR	<table border="1"><tr><td>0111</td><td>SR</td><td>BaseR</td><td colspan="3">offset6</td></tr></table>	0111	SR	BaseR	offset6			STR SR, BaseR, offset6
1111	0000	trapvect8															
0111	SR	BaseR	offset6														
		$R7 \leftarrow PC, PC \leftarrow M[ZEXT(trapvect8)]$			$M[BaseR + SEXT(offset6)] \leftarrow SR$												

End of ECE 220 Final Exam